

Satisficing Models to Mitigate Uncertainty

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outline

- optimization under uncertainty and satisficing - a quick review
- satisficing decision criteria - general representation theorem
- the t-model: a tractable probabilistic satisficing model
- numerical illustration - maximum coverage facility location problem

optimization under uncertainty

- deterministic optimization:

$$\min_{\mathbf{x} \in \mathcal{X}} \mathbf{c}'\mathbf{x} \quad \text{s.t.} \quad \mathbf{A}\mathbf{x} \geq \mathbf{b}$$

- optimization under uncertainty $\tilde{\mathbf{z}} \in \mathcal{W}$:

$$\min_{\mathbf{x} \in \mathcal{X}} \mathbf{c}'\mathbf{x} \quad \text{s.t.} \quad \mathbf{A}(\tilde{\mathbf{z}})\mathbf{x} \geq \mathbf{b}(\tilde{\mathbf{z}})$$

- probability distributions available: stochastic optimization (Prékopa 1995; Birge and Louveaux 1997; etc.)
- distributions unavailable:
 - robust optimization (Ben-Tal and Nemirovski 1999; Bertsimas and Sim 2004, Bertsimas et al. 2011)
 - distributionally robust optimization (Delage and Ye 2010; Wiesemann et al. 2014; etc ... this workshop)

satisficing

satisficing = satisfy + suffice

Simon (1959):

«... the entrepreneur may not care to maximize, but may simply want to earn a return that he regards as satisfactory [...] “satisfactory profits” is a concept more meaningfully related to the psychological notion of aspiration levels than to maximization...»

Simon, H. A. (1959). Theories of Decision-Making in Economics and Behavioral Science. *The American Economic Review* 49(3):253–83.

a first satisficing model under uncertainty: the p-model

Charnes and Cooper (1963)

- first to incorporate the idea of satisficing in mathematical programming under uncertainty:

$$\begin{aligned} \max \quad & \ln \mathbb{P}(\mathbf{A}(\tilde{\mathbf{z}})\mathbf{x} \geq \mathbf{b}(\tilde{\mathbf{z}})) \\ \text{s.t.} \quad & \mathbf{x} \in \mathcal{X}. \end{aligned}$$

- randomly perturbed linear constraints $\mathbf{A}(\tilde{\mathbf{z}})\mathbf{x} \geq \mathbf{b}(\tilde{\mathbf{z}})$
 - tractable only for restricted special cases
 - general case intractable (Nemirovski and Shapiro 2006)

Charnes, A., and W. Cooper (1963) Deterministic Equivalents for Optimizing and Satisficing under Chance Constraints. *Operations Research* 11(1):18–39.

chance-constrained optimization/programming

Charnes and Cooper (1959), close relation to the p-model

$$\begin{aligned} \min \quad & \mathbf{c}'\mathbf{x} \\ \text{s.t.} \quad & \ln \mathbb{P}(\mathbf{A}(\tilde{\mathbf{z}})\mathbf{x} \geq \mathbf{b}(\tilde{\mathbf{z}})) \geq \Delta \\ & \mathbf{x} \in \mathcal{X} \end{aligned}$$

- satisficing criterion subject to a lower bound parameter $\Delta \in \mathbb{R}$
- objective is a deterministic cost function; $\mathbf{c} \in \mathbb{R}^N$ defines the objective function coefficients
- approximation by sample average approximation (SAA) methods
 - disadvantage: require large number of samples

Charnes, A., and W. Cooper (1959) Chance-Constrained Programming. Management Science 6(1):73–79.

robust optimization

$$\begin{aligned}
 \min \quad & \mathbf{c}'\mathbf{x} \\
 \text{s.t.} \quad & \mathbf{A}(\mathbf{z})\mathbf{x} \geq \mathbf{b}(\mathbf{z}) \quad \forall \mathbf{z} \in \mathcal{U}(\Gamma) \\
 & \mathbf{x} \in \mathcal{X}
 \end{aligned}$$

- \mathbf{z} denotes realization of $\tilde{\mathbf{z}}$ from an *uncertainty set*, $\mathcal{U}(\Gamma)$ (typically $\mathcal{U}(\cdot)$ designed such that $\mathcal{U}(\alpha_1) \subseteq \mathcal{U}(\alpha_2) \subseteq \mathcal{W} \subseteq \mathbb{R}^K$ for all $0 \leq \alpha_1 \leq \alpha_2$).
- does not require the specification of a probability distribution, but instead a “budget of uncertainty” $\Gamma \in \mathbb{R}_+$
 - the level of uncertainty that must be tolerated
 - may not be easy to specify
- yields tractable formulations under reasonable conditions: e.g., if $\mathcal{U}(\Gamma)$ is described as norm-based sets $\mathcal{U}(\Gamma) = \{\mathbf{z} \in \mathcal{W} \mid \|\mathbf{z}\| \leq \Gamma\}$:
 - linear program for $\|\cdot\|_1$, $\|\cdot\|_\infty$ and D-Norm (Bertsimas et al. 2004)
 - second-order cone program for $\|\cdot\|_2$ norm

a satisficing model for robust optimization

- the **p-model** is a satisficing model for chance-constrained optimization:

$$\begin{aligned} \max \quad & \ln \mathbb{P}(\mathbf{A}(\tilde{\mathbf{z}})\mathbf{x} \geq \mathbf{b}(\tilde{\mathbf{z}})) \\ \text{s.t.} \quad & \mathbf{x} \in \mathcal{X} \end{aligned}$$

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- can we define a satisficing model for robust optimization?

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- can we define a satisficing model for robust optimization?

→ the **r-model**:

$$\begin{aligned} \max \quad & \{\alpha \mid \mathbf{A}(\mathbf{z})\mathbf{x} \geq \mathbf{b}(\mathbf{z}) \quad \forall \mathbf{z} \in \mathcal{U}(\alpha)\} \\ \text{s.t.} \quad & \mathbf{x} \in \mathcal{X} \end{aligned}$$

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some benefits of satisficing models

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$$\begin{aligned} \min \quad & \mathbf{c}'\mathbf{x} \\ \text{s.t.} \quad & \ln \mathbb{P}(\mathbf{A}(\tilde{\mathbf{z}})\mathbf{x} \geq \mathbf{b}(\tilde{\mathbf{z}})) \geq \Delta \\ & \mathbf{x} \in \mathcal{X} \end{aligned}$$

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satisficing decision criterion - definition

setting:

- $\tilde{\mathbf{z}}$ a K dimensional random vector that influences the entries of the function maps $\mathbf{A} : \mathbb{R}^K \mapsto \mathbb{R}^{M \times N}$ and $\mathbf{b} : \mathbb{R}^K \mapsto \mathbb{R}^M$.
- randomly perturbed linear constraints, $\mathbf{A}(\mathbf{z})\mathbf{x} \geq \mathbf{b}(\mathbf{z})$, where \mathbf{z} is a random outcome of $\tilde{\mathbf{z}}$.
- $\mathcal{W} \subseteq \mathbb{R}^K$ the support of the random vector $\tilde{\mathbf{z}}$.

definition: satisficing decision criterion

a function $\nu : \mathbb{R}^N \mapsto \mathbb{R} \cup \{-\infty\}$ is a *satisficing decision criterion* if it has the following two properties. For all $\mathbf{x}, \mathbf{y} \in \mathbb{R}^N$,

- 1 (satisficing dominance) if $\mathbf{A}(\mathbf{z})\mathbf{y} \geq \mathbf{b}(\mathbf{z})$ implies $\mathbf{A}(\mathbf{z})\mathbf{x} \geq \mathbf{b}(\mathbf{z})$ for all $\mathbf{z} \in \mathcal{W}$, then $\nu(\mathbf{x}) \geq \nu(\mathbf{y})$.
- 2 (infeasibility) if there does not exist $\mathbf{z} \in \mathcal{W}$ such that $\mathbf{A}(\mathbf{z})\mathbf{x} \geq \mathbf{b}(\mathbf{z})$, then $\nu(\mathbf{x}) = -\infty$.

satisficing decision criteria - previous examples

- the p-model is an optimization problem that maximizes a satisficing decision criterion $\nu_P : \mathbb{R}^N \mapsto \mathbb{R} \cup \{-\infty\}$ given by

$$\nu_P(\mathbf{x}) = \ln \mathbb{P}(\mathbf{A}(\tilde{\mathbf{z}})\mathbf{x} \geq \mathbf{b}(\tilde{\mathbf{z}}))$$

- the r-model is an optimization problem that maximizes a satisficing decision criterion $\nu_R : \mathbb{R}^N \mapsto \mathbb{R} \cup \{-\infty\}$ given by

$$\nu_R(\mathbf{x}) = \max \{ \alpha \mid \mathbf{A}(\mathbf{z})\mathbf{x} \geq \mathbf{b}(\mathbf{z}) \quad \forall \mathbf{z} \in \mathcal{U}(\alpha) \}$$

satisficing decision criteria - a general representation

a general representation of any satisficing decision criterion ν can be given by the following result:

theorem: general representation

consider a function $\nu : \mathbb{R}^N \mapsto \mathbb{R} \cup \{-\infty\}$ defined as

$$\nu(\mathbf{x}) = \max_{\alpha \in \mathcal{S}} \{ \rho(\alpha) \mid \mathbf{A}(\mathbf{z})\mathbf{x} \geq \mathbf{b}(\mathbf{z}) \quad \forall \mathbf{z} \in \mathcal{U}(\alpha) \} \quad (1)$$

for some function $\rho : \mathcal{S} \rightarrow \mathbb{R} \cup \{-\infty\}$ on domain $\mathcal{S} \subseteq \mathbb{R}^P$, and for some family of nonempty uncertainty sets $\mathcal{U}(\alpha) \subseteq \mathcal{W}$ defined for all $\alpha \in \mathcal{S}$; then the function ν is a satisficing decision criterion; moreover, any satisficing decision criterion can be represented in a form given by (1) with $\mathcal{S} \subseteq \mathbb{R}^N$.

the s-model: a general family of satisficing models

general s-model

$$\begin{aligned}
 \max \quad & \rho(\alpha) \\
 \text{s.t.} \quad & \mathbf{A}(\mathbf{z})\mathbf{x} \geq b(\mathbf{z}) \quad \forall \mathbf{z} \in \mathcal{U}(\alpha) \\
 & \mathbf{x} \in \mathcal{X} \\
 & \alpha \in \mathcal{S}
 \end{aligned}$$

- adjusts uncertainty sets $\mathcal{U}(\alpha)$ for which the constraints remain feasible
- maximizes $\rho(\alpha) : \mathcal{S} \rightarrow \mathbb{R}$
- careful design of $\rho(\alpha)$ and $\mathcal{U}(\alpha)$ can lead to meaningful and tractable models

the probabilistic s-model

- recap - the most general satisficing model:

general s-model

$$\begin{aligned}
 \max \quad & \rho(\alpha) \\
 \text{s.t.} \quad & \mathbf{A}(\mathbf{z})\mathbf{x} \geq b(\mathbf{z}) \quad \forall \mathbf{z} \in \mathcal{U}(\alpha) \\
 & \mathbf{x} \in \mathcal{X} \\
 & \alpha \in \mathcal{S}
 \end{aligned}$$

- how to combine useful aspects of both the p-model and the r-model?
- set $\rho(\alpha) = \ln \mathbb{P}(\tilde{\mathbf{z}} \in \mathcal{U}(\alpha))$

probabilistic s-model

$$\begin{aligned}
 \max \quad & \ln \mathbb{P}(\tilde{\mathbf{z}} \in \mathcal{U}(\alpha)) \\
 \text{s.t.} \quad & \mathbf{A}(\mathbf{z})\mathbf{x} \geq b(\mathbf{z}) \quad \forall \mathbf{z} \in \mathcal{U}(\alpha) \\
 & \mathbf{x} \in \mathcal{X} \\
 & \alpha \in \mathcal{S}
 \end{aligned}$$

the t-model: a tractable probabilistic s-model

- uncertain parameters \tilde{z}_k , $k \in [K]$ are independently distributed real random variables with support \mathcal{W}_k ; $\mathcal{W} = \times_{k=1}^K \mathcal{W}_k$
- uncertainty defined by affine functions:

$$\mathbf{a}_i(\mathbf{z}) = \mathbf{a}_i^0 + \sum_{k=1}^K \mathbf{a}_i^k z_k \text{ and } b_i(\mathbf{z}) = b_i^0 + \sum_{k=1}^K b_i^k z_k$$
- family of adjustable uncertainty sets (“box” type):

$$\mathcal{U}(\boldsymbol{\alpha}) = \mathcal{U}(\underline{\boldsymbol{\alpha}}, \bar{\boldsymbol{\alpha}}) = \left\{ \mathbf{z} \in \mathbb{R}^K : \mathbf{z} \in [\underline{\boldsymbol{\alpha}}, \bar{\boldsymbol{\alpha}}] \right\}.$$

t-model

$$\begin{aligned} \max \quad & \sum_{k \in [K]} \ln \mathbb{P}(\underline{\alpha}_k \leq \tilde{z}_k \leq \bar{\alpha}_k) \\ \text{s.t.} \quad & \mathbf{A}(\mathbf{z})\mathbf{x} \geq \mathbf{b}(\mathbf{z}) \quad \forall \mathbf{z} \in [\underline{\boldsymbol{\alpha}}, \bar{\boldsymbol{\alpha}}] \\ & \mathbf{x} \in \mathcal{X}, \quad \underline{\boldsymbol{\alpha}} \leq \bar{\boldsymbol{\alpha}}, \quad \underline{\boldsymbol{\alpha}}, \bar{\boldsymbol{\alpha}} \in \mathcal{W} \end{aligned}$$

reformulation (*robust counterpart*) of the t-model

reformulation: t-model

$$\begin{aligned}
 \max \quad & \sum_{k \in [K]} \ln \mathbb{P}(\underline{\alpha}_k \leq \tilde{\mathbf{z}} \leq \bar{\alpha}_k) \\
 \text{s.t.} \quad & \sum_{j \in [M]} a_{ij}^0 x_j + \sum_{k \in [K]} v_{ik} \geq b_i^0 \quad \forall i \in [M] \\
 & v_{ik} \leq \sum_{j \in [M]} a_{ij}^k x_j \bar{\alpha}_k - b_i^k \bar{\alpha}_k \quad \forall i \in [M], k \in [K] \\
 & v_{ik} \leq \sum_{j \in [M]} a_{ij}^k x_j \underline{\alpha}_k - b_i^k \underline{\alpha}_k \quad \forall i \in [M], k \in [K] \\
 & \mathbf{x} \in \mathcal{X}, \mathbf{v} \in \mathbb{R}^{M \times K}, \\
 & \underline{\alpha} \leq \bar{\alpha}, \underline{\alpha}, \bar{\alpha} \in \mathcal{W}.
 \end{aligned}$$

- polynomial number of constraints (good)
- remaining difficulties:
 - non-linear objective function.
 - the terms $x_j \underline{\alpha}_k$ and $x_j \bar{\alpha}_k$ $j \in [M], k \in [K]$ are bilinear.

t-model for log-concave densities

the t-model is tractable if:

- the distributions of the random variables are described by log-concave densities
- constraints are linear: e.g., uncertainty in right-hand-side only or, decision variables \mathbf{x} are binary

note: consequence for the non-linear objective function:

- if \tilde{z}_k is log-concave, then $\ln \mathbb{P}(\underline{\delta} \leq \tilde{z}_k \leq \bar{\delta})$ is a concave function of $(\underline{\delta}, \bar{\delta})$
- the objective function can be approximated by piecewise linear approximation of arbitrary accuracy (density cuts) \rightarrow branch-and-cut

t-model for discrete distributions

- $\mathcal{W}_k = \{\zeta_k^1, \zeta_k^2, \dots, \zeta_k^{L(k)}\}$, $\mathbb{P}(\tilde{z}_k = \zeta_k^\ell) = p_k^\ell$
- outcomes ζ_k^ℓ sorted in non-decreasing order
- $\mathcal{U}(\alpha) = \left\{ \mathbf{z} \in \mathcal{W} \mid \sum_{\ell \in [L(k)]} \zeta_k^\ell \underline{\alpha}_k^\ell \leq z_k \leq \sum_{\ell \in [L(k)]} \zeta_k^\ell \bar{\alpha}_k^\ell, \forall k \in [K] \right\}$

t-model for discrete distributions

$$\begin{aligned}
 & \max \quad \ln \mathbb{P}(\tilde{\mathbf{z}} \in \mathcal{U}(\alpha)) \\
 & \text{s.t.} \quad \mathbf{A}(\mathbf{z})\mathbf{x} \geq \mathbf{b}(\mathbf{z}) \quad \forall \mathbf{z} \in \mathcal{U}(\alpha) \\
 & \quad \sum_{\ell \in [L(k)]} \underline{\alpha}_k^\ell = 1, \quad \sum_{\ell \in [L(k)]} \bar{\alpha}_k^\ell = 1, \quad \forall k \in [K] \\
 & \quad \sum_{\ell \in [L(k)]} \ell(\bar{\alpha}_k^\ell - \underline{\alpha}_k^\ell) \geq 0, \quad \forall k \in [K] \\
 & \quad \underline{\alpha}_k, \bar{\alpha}_k \in \{0, 1\}^{L(k)} \quad \forall k \in [K], \quad \mathbf{x} \in \mathcal{X}.
 \end{aligned}$$

monotone t-models

a t-model significantly simplifies if it is *monotone*:

definition

a t-model is *monotone* with respect to the uncertain parameters $\tilde{\mathbf{z}}$ if there exists a partition $\overline{\mathcal{K}}, \underline{\mathcal{K}} \subseteq [K]$, i.e., $\overline{\mathcal{K}} \cap \underline{\mathcal{K}} = \emptyset$, $\overline{\mathcal{K}} \cup \underline{\mathcal{K}} = [K]$ such that for all $k \in \overline{\mathcal{K}}$

$$\sum_{j \in [M]} a_{ij}^k x_j \leq b_i^k \quad \forall i \in [M], \mathbf{x} \in \mathcal{X}$$

and for all $k \in \underline{\mathcal{K}}$

$$\sum_{j \in [M]} a_{ij}^k x_j \geq b_i^k \quad \forall i \in [M], \mathbf{x} \in \mathcal{X}$$

a monotone t-model can also be turned into an adjustable t-model for multi-stage decision making !!

adjustable t-model for multi-stage decision making

- $(T + 1)$ -stage problem.
- first stage, decision $\mathbf{x}^0 \in \mathbb{R}^{N_0}$ is made before any uncertainty is realized.
- In subsequent stages, decisions made are $\mathbf{x}^1(\tilde{\mathbf{z}}_{\mathcal{T}_1}), \dots, \mathbf{x}^T(\tilde{\mathbf{z}}_{\mathcal{T}_T})$, where the *recourse decision* \mathbf{x}^t at stage $t + 1$ is a measurable function $\mathbf{x}^t : \mathbb{R}^{|\mathcal{T}_t|} \mapsto \mathbb{R}^{N_t}$ that maps from the realization of the uncertain parameters $\tilde{\mathbf{z}}_{\mathcal{T}_t}$ to the appropriate action in \mathbb{R}^{N_t} .
- let

$$\mathbf{A}(\mathbf{z}) = [\mathbf{A}^0(\mathbf{z}) \ \mathbf{A}^1(\mathbf{z}) \ \dots \ \mathbf{A}^T(\mathbf{z})], \quad \mathbf{x}(\mathbf{z}) = (\mathbf{x}^0, \mathbf{x}^1(\mathbf{z}_{\mathcal{T}_1}), \dots, \mathbf{x}^T(\mathbf{z}_{\mathcal{T}_T}))$$

of appropriate dimensions so that

$$\mathbf{A}(\mathbf{z})\mathbf{x}(\mathbf{z}) = \mathbf{A}^0(\mathbf{z})\mathbf{x}^0 + \sum_{t \in [T]} \mathbf{A}^t(\mathbf{z})\mathbf{x}^t(\mathbf{z}_{\mathcal{T}_t}).$$

adjustable t-model for multi-stage decision making, cont.

- formulate the adjustable T-model as follows:

$$\begin{aligned}
 \max \quad & \sum_{k \in [K]} \ln \mathbb{P}(\underline{\alpha}_k \leq \tilde{z}_k \leq \bar{\alpha}_k) \\
 \text{s.t.} \quad & \mathbf{A}(\mathbf{z})\mathbf{x}(\mathbf{z}) \geq \mathbf{b}(\mathbf{z}) && \forall \mathbf{z} \in [\underline{\alpha}, \bar{\alpha}] \\
 & \mathbf{x}(\mathbf{z}) \in \mathcal{X} && \forall \mathbf{z} \in \mathcal{W} \\
 & \mathbf{x}^t \in \mathcal{R}(|\mathcal{T}_t|, N_t) && \forall t \in [T] \\
 & \underline{\alpha} \leq \bar{\alpha}, \underline{\alpha}, \bar{\alpha} \in \mathcal{W},
 \end{aligned}$$

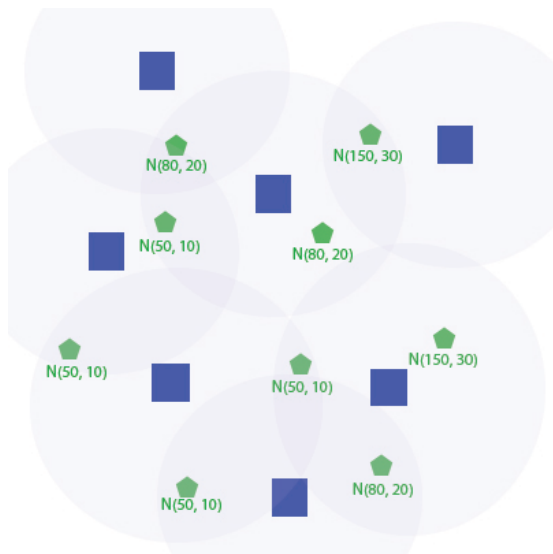
where $\mathcal{R}(m, n)$ denotes the family of all measurable functions that map from \mathbb{R}^m to \mathbb{R}^n .

- under fixed recourse assumptions and our monotonicity condition, equivalent to solving a one-stage problem.

stochastic maximum coverage facility location problem

given:

- \mathcal{I} candidate facility locations; \mathcal{J} customer demands \tilde{d}_j
- network is “sparse”: each customer can be covered by approx. 15% - 40% of all facilities
- available budget B ; facility construction costs c_j ; capacities a_j



stochastic maximum coverage facility location problem

select facilities that maximize the probability that all demands can be satisfied:

initial p-model for maximum coverage problem, hard to solve

$$\begin{array}{ll}
 \max & \ln \mathbb{P} \left(\sum_{i \in \mathcal{I}_j} y_{ij}(\tilde{\mathbf{z}}) \geq \tilde{z}_j \quad \forall j \in \mathcal{J} \right) \\
 \text{s.t.} & \sum_{j \in \mathcal{J}_i} y_{ij}(\mathbf{z}) \leq a_i x_i \quad \forall \mathbf{z} \in \mathcal{W}, i \in \mathcal{I} \\
 & \sum_{i \in \mathcal{I}} c_i x_i \leq B \\
 & y_{ij}(\mathbf{z}) \geq 0 \quad \forall \mathbf{z} \in \mathcal{W}, i \in \mathcal{I}, j \in \mathcal{J}_i \\
 & y_{ij}(\cdot) \in \mathcal{R}(|\mathcal{J}|, 1) \quad \forall i \in \mathcal{I}, j \in \mathcal{J}_i \\
 & x_i \in \{0, 1\} \quad \forall i \in \mathcal{I},
 \end{array}$$

t-models and monte carlo benchmarks

t-models:

- T-1: branch-and-cut for log-concave densities
 - maximizes the probability that each demand is met
 - assumes knowledge of the probability distribution
- T-2: sample based model (discrete distribution)
 - L data samples (scenarios)
 - maximizes # of outcomes that are feasible in constraints
 - no assumptions about probability distributions

SAA models with L data samples:

- P-1: maximizing feasibility probability
 - maximizes the number of feasible scenarios (obj.: P-model)
- E: minimizing expected demand shortfall
 - minimizes the expected demand shortfall (obj.: expected value)

computational settings

problem instances (total of 60):

- # customers $|\mathcal{J}| \in \{100, 250, 500, 1000, 2000\}$
- # facilities $|\mathcal{I}| \in \{0.5|\mathcal{J}|, |\mathcal{J}|, 2|\mathcal{J}|\}$
- network density $A_p \in \{15, 20, 30, 40\}$
- $\tilde{d}_j \sim N(\mu_j, (0.5\mu_j)^2)$; $\mu_j \sim U(1, 100)$
- budget B set 1.05 times the costs required to satisfy the average demand

computational settings:

- CPLEX 12.6.1 with standard parameters
- 12hrs computing time limit, 24gb memory limit
- evaluation via Monte Carlo simulation (100,000 samples)

customers - scalability of the model T-1

$ \mathcal{J} $	succ. rate %	demand shortfall	time (minutes)
100	84.41	2.0	22.7
250	82.76	4.2	2.7
500	99.93	0.0	3.4
1000	96.60	1.2	30.7
2000	95.97	2.2	426.9
all	92.06	1.9	98.5

Table: Out-of-sample performance study for different problem sizes, reporting average success rate (%), average demand shortfall (in 10 units), average computing time (in minutes)

- solves all instances
- high success rates and low shortfalls for all problem sizes
- reasonable computing times

L - scalability: data samples based models

average over all 60 instances

L	succ. rate %	T-2 short fall	time	# ns	succ. rate %	P-1 short fall	time	# ns	succ. rate %	E short fall	time	# ns
5	88.73	350.9	44.9	2	64.65	2,440.8	207.0	16	70.10	1,743.6	181.2	12
10	85.39	699.8	62.5	4	43.99	3,335.8	379.1	29	52.69	2,826.7	354.8	22
15	88.41	351.0	75.3	2	35.06	3,523.0	448.7	35	41.85	3,037.6	463.2	30
50	86.64	525.4	54.9	3	20.37	3,627.7	606.7	45	26.54	3,370.2	569.8	39

- # ns: number of instances without feasible solution
- models P-1 and E hard to solve as L increases
- model T-2 remains relatively stable

L - scalability & robustness

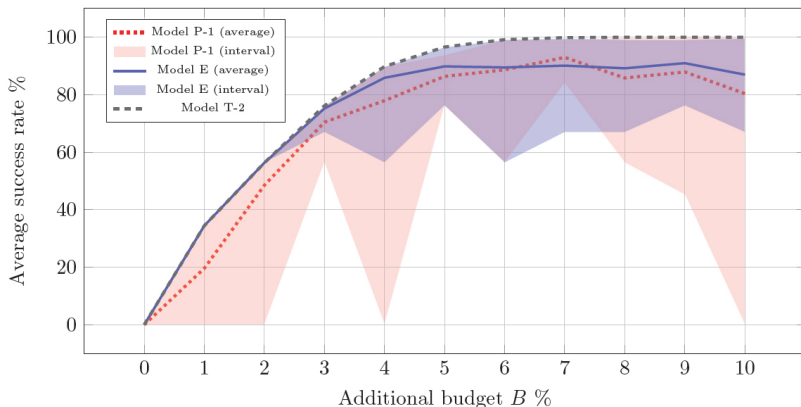
average [min, max] among averages of 10 replications increasing L: all instances

J	P-1	E	T-2				
	L=5	L=5	L=5	L=15	L=50	L=100	L=500
100	80.70 [71.7, 86.8]	76.29 [66.4, 85.8]	80.41 [64.0, 87.3]	80.56 [70.7, 86.3]	84.69 [82.2, 86.9]	84.57 [81.3, 87.1]	84.47 [80.6, 86.8]
250	81.11 [69.1, 82.8]	82.51 [82.3, 82.8]	80.44 [75.4, 82.7]	80.80 [75.6, 82.5]	81.57 [75.8, 82.7]	82.38 [81.9, 82.5]	82.43 [81.8, 82.5]
500	96.66 [91.7, 99.9]	95.69 [86.9, 99.9]	98.29 [91.8, 99.9]	99.92 [99.9, 99.9]	99.92 [99.9, 99.9]	99.92 [99.9, 99.9]	99.92 [99.9, 99.9]
1000	47.90 [23.7, 63.4]	70.96 [60.2, 88.5]	96.58 [96.5, 96.7]	96.62 [96.6, 96.6]	96.63 [96.6, 96.6]	96.63 [96.6, 96.6]	96.63 [96.6, 96.6]
2000	0.77 [0.0, 7.7]	26.84 [20.3, 31.3]	73.97 [63.0, 86.5]	72.41 [63.0, 78.7]	73.19 [63.0, 86.6]	82.63 [63.0, 94.5]	77.93 [63.0, 94.5]
all	61.25 [56.1, 65.2]	70.40 [67.6, 77.4]	86.03 [82.5, 88.8]	86.15 [84.0, 88.5]	87.24 [85.4, 88.9]	89.30 [85.8, 91.7]	88.34 [85.5, 91.9]

Table: Comparison of average [minimum, maximum] success rates (%) over all problem instances among 10 replications for Models P-1 and E with same sample size $L = 5$, and for Model T-2 with different sample sizes.

investment study: success rates

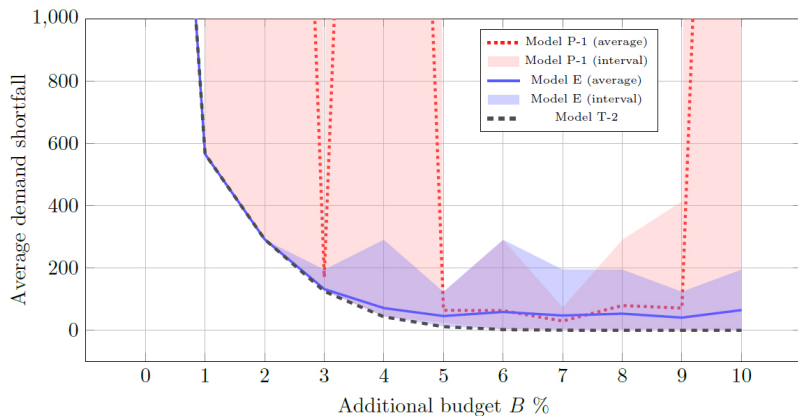
average intervals between min/max among 10 runs
 ($L = 5$, $|\mathcal{I}| = 500$, $|\mathcal{J}| = 1000$)



- P-1 and E models unstable; T-2 model stable at highest success rates

investment study: demand shortfall

average intervals between min/max among 10 runs
 ($L = 5$, $|\mathcal{I}| = 500$, $|\mathcal{J}| = 1000$)



- P-1 and E models unstable; T-2 model stable at lowest shortfalls

conclusions

contributions:

- introduction of the s-model
 - flexible adjustment of uncertainty sets
 - generalizes the p-model
 - provides link to general chance-constraints and robust optimization problems
- general framework: allows for many tractable implementations, e.g., the t-model
- exemplified for continuous and discrete/empiric distributions
 - log-concave density functions: cut-based solution methods
 - data sampling/discrete distributions: efficient reformulation to mixed-integer program

conclusions

contributions - computational experiments:

- maximum coverage facility location problem
- large problem instances
- benchmark approaches (SAA) cannot handle large sample sizes
- t-models scale well for all instances
 - knowledge about probability distribution helps
 - without available distributions, large sample size results in stable results for all instances

future research directions

- simple idea and easy to implement
 - relevant for decision makers in practice
 - applicable to many (difficult) problems
 - high performance → competitive alternative to traditional sampling methods
- cut-based method has been explored for NP-hard MIP
 - likely to be very quick for linear programs (cutting plane)
- further implementations of the S-model
 - other cases may yield tractable models for important problems
- scalability of data-driven approach
 - may handle even larger data sets when solved by advanced optimization methods - big data/machine learning?